

Extremal set theory

Let A be a base set. Then 2^A is the power set of A , so it contains all subsets of A .
 $\mathcal{F} \subseteq 2^A$ is called a set-system or a family of subsets.

We do not lose generality if we say that $A = \{1, 2, 3, \dots, n\} = [n]$.

A set-system can be interpreted as a hypergraph where the vertices are the elements of A , the hyperedges are the elements of \mathcal{F} , which are subsets of A .

A set-system/hypergraph is k -uniform if each element of \mathcal{F} /hyperedge is a k element set.
 A 2-uniform set-system/hyperedge is a graph.

Extremal set theory: How big can \mathcal{F} be if it satisfies some conditions?

Examples:

$\mathcal{F} \subseteq 2^{[n]}$, for any pair $A, B \in \mathcal{F}$: if $A \neq B \Rightarrow A \cap B = \emptyset \Rightarrow \max |\mathcal{F}| = n+1$

The maximum is attained when \mathcal{F} contains each one element subset and the empty set.

$\mathcal{F} \subseteq 2^{[n]}$, for any pair $A, B \in \mathcal{F}$: if $A \neq B \Rightarrow A \cap B \neq \emptyset \Rightarrow \max |\mathcal{F}| = 2^{n-1}$

Proof: If $A \in \mathcal{F}$ then $\overline{A} \notin \mathcal{F}$, therefore at most half of the subsets of $[n]$ can be included in \mathcal{F} . And this bound can be attained if we choose all subsets of $[n]$ which contain 1.

Theorem: Erdős-Ko-Rado 1961: $\mathcal{F} \subseteq 2^{[n]}$ and \mathcal{F} is k -uniform where $k \leq n/2$ and for any pair $A, B \in \mathcal{F}$: $A \cap B \neq \emptyset$ then $|\mathcal{F}| \leq \binom{n-1}{k-1}$.

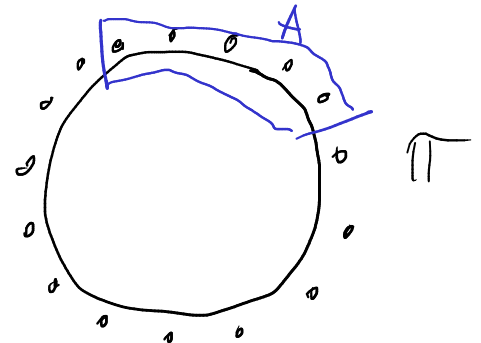
Note that if $k > n/2$, then any two k element subsets of $[n]$ intersect each other, this is why the opposite case is interesting.

$\binom{n-1}{k-1}$ can be attained if we pick all the k element subsets of $[n]$ which contain 1.

Proof by Gyula Katona OH '72:

Let π be a circular permutation of $[n]$. So we imagine that the n elements are seated around a table.

A set A is an arc in π if its elements are consecutive in π :



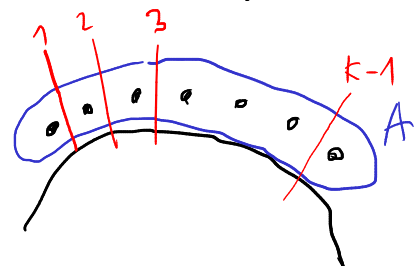
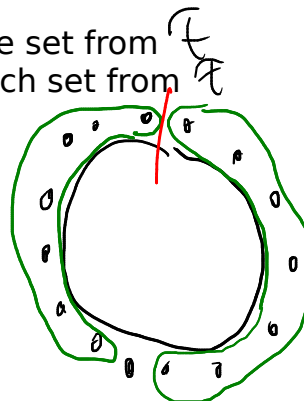
Fix a circular permutation π .

There are at most k sets from \mathcal{F} which are arcs in π :

Pick a set A from \mathcal{F} which is an arc in π . An other set B from \mathcal{F} which is also an arc in π can end at $k-1$ different positions, because B intersect A .

For each such position there is only one set from \mathcal{F} which is an arc in π and ends there: Each set from \mathcal{F} has the same size k , therefore two arcs B and C of size k which end at the same place are either the same or disjoint because $k \leq n$.

So for π there are at most k set from \mathcal{F} which are arcs.



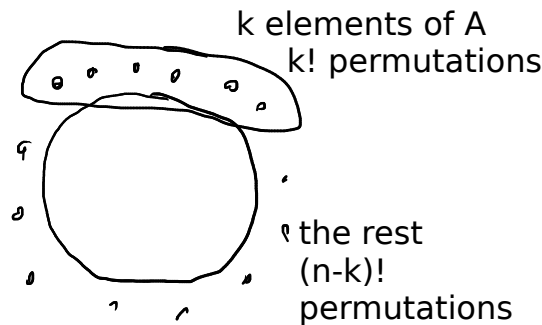
The total number of circular permutations is $(n-1)!$

So if we consider all of the circular permutations, we find at most $k(n-1)!$ arcs which correspond to sets from \mathcal{F} .

A set from \mathcal{F} is an arc in $k!(n-k)!$ circular permutations. Therefore:

$$|\mathcal{F}| \leq \frac{k(n-1)!}{k!(n-k)!} = \binom{n-1}{k-1}$$

□



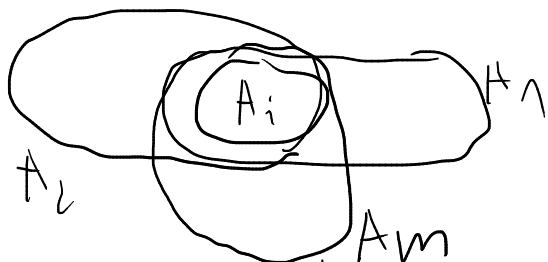
Theorem: Fischer's inequality:

$$\mathcal{F} = \{A_1, A_2, \dots, A_m\} \subseteq 2^U, \lambda > 0 \text{ and for } \forall i \neq j: |A_i \cap A_j| = \lambda \Rightarrow m \leq n$$

Proof: For each i $|A_i| \geq \lambda$.

Case 1: There is an element A_i s.t. $|A_i| = \lambda$. Then every other set A_j contains A_i as a subset. But then $\{A_1 \setminus A_i, A_2 \setminus A_i, \dots, A_{i-1} \setminus A_i, A_{i+1} \setminus A_i, \dots, A_m \setminus A_i\}$

is an m element set-system such that any two elements are disjoint, and its base set contains $n - \lambda + 1 \leq n$ sets, so $m \leq n$.



Case 2: $\forall i: |A_i| > \lambda$:

Let \bar{a}_i be the characteristic vector of A_i , so it is an n dimensional vector and $a_{i,j} = \begin{cases} 1 & \text{if } j \in A_i \\ 0 & \text{if } j \notin A_i \end{cases}$

Example: If $n=7$ and $A_i = \{1, 3, 5\}$ then $\bar{a}_i = [1, 0, 1, 0, 1, 0, 0]$

We are going to show that $\bar{a}_1, \bar{a}_2, \dots, \bar{a}_m$ are linearly independent.

$$\bar{a}_i \cdot \bar{a}_j = \begin{cases} \lambda & \text{if } i \neq j \\ |A_i| & \text{if } i = j \end{cases}$$

Consider a linear combination $c_1 \bar{a}_1 + c_2 \bar{a}_2 + c_3 \bar{a}_3 + \dots + c_m \bar{a}_m = \mathbf{0}$

Take the dot product of this vector and \bar{a}_i :

$$0 = \mathbf{0} \cdot \bar{a}_i = (c_1 \bar{a}_1 + c_2 \bar{a}_2 + \dots + c_m \bar{a}_m) \cdot \bar{a}_i = (c_1 + c_2 + \dots + c_m) \lambda + c_i (|A_i| - \lambda)$$

If $c_1 + c_2 + c_3 + \dots + c_m = 0 \Rightarrow c_i = 0 \forall i$

If $c_1 + c_2 + c_3 + \dots + c_m > 0 \Rightarrow c_i < 0 \forall i$

If $c_1 + c_2 + c_3 + \dots + c_m < 0 \Rightarrow c_i > 0 \forall i$

$\bar{a}_1, \bar{a}_2, \dots, \bar{a}_m$ are linearly independent

If m vectors which are n dimensional are linearly independent, then $m \leq n$.

□

Theorem: Ray-Chaudhuri-Wilson:

$$L = \{l_1, \dots, l_s\}, \mathcal{F} \subseteq 2^{[n]}, \forall A, B \in \mathcal{F}, A \neq B: |A \cap B| \in L \text{ then: } |\mathcal{F}| \leq \sum_{i=0}^s \binom{n}{i}$$

This bound can be attained, for example when $L = \{0, 1, 2, \dots, s-1\}$, then $\hat{\mathcal{F}} = \{A \subseteq [n] \mid |A| \leq s\}$ and $|\hat{\mathcal{F}}| = \sum_{i=0}^s \binom{n}{i}$.

When $s=1$, Ray-Chaudhuri-Wilson gives that $\forall A, B \in \mathcal{F}, A \neq B: |A \cap B| = 0 \Rightarrow |\mathcal{F}| \leq \sum_{i=0}^1 \binom{n}{i} = n+1$

Fischer's inequality gives: $\forall A, B \in \mathcal{F}, A \neq B: |A \cap B| = \lambda \Rightarrow |\mathcal{F}| \leq n$

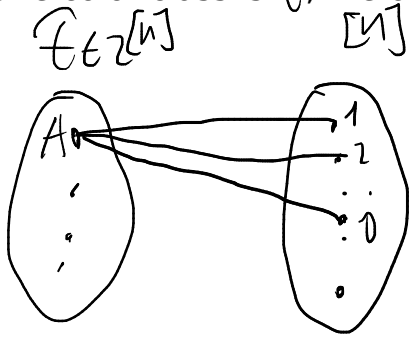
The difference comes from that $\lambda \geq 0$ but $\lambda > 0!$

Theorem: De Bruijn-Erdős: $\mathcal{G} \subseteq 2^{[m]}, \forall i, j \in [m], i \neq j$: there is exactly one set $A \in \mathcal{G}$ such that $i, j \in A$. and if $B \in \mathcal{G} \Rightarrow |B| \geq 2$.
 $\Rightarrow |\mathcal{G}| = 1$ or $|\mathcal{G}| \geq m$.

This theorem can be interpreted in the following way: If m points in the projective plane are given then either they are on the same line or they determine at least m different lines.

We prove this theorem by using hypergraph duality:

Any hypergraph $\mathcal{F} \subseteq 2^{[n]}$ can be represented by a bipartite graph in the following way: One color class is \mathcal{F} , the other one is $[n]$, and a set A is incident to an element j if $j \in A$.

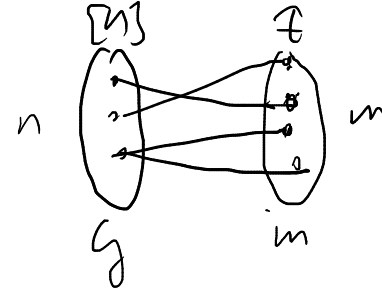


this can be a multiset of sets

if $\mathcal{F} \subseteq 2^{[n]}, |\mathcal{F}| = m$ then the dual hypergraph of \mathcal{F} is a hypergraph $\mathcal{G} \subseteq 2^{[m]}, 1 \leq |\mathcal{G}| \leq n$.

$|\mathcal{G}| = n$ and the bipartite graph representation of \mathcal{G} and \mathcal{F} are the same, except

that the role of the two color classes have been swapped. So the class which contained elements of \mathcal{F} is containing the elements of $[m]$ and the other class which contained the elements of $[n]$ is containing the elements of \mathcal{G} .



Example:
 $\mathcal{F} \subseteq 2^{\{1,2,3\}}, \mathcal{F} = \{A, B, C\}$ $\mathcal{G} \subseteq 2^{\{A,B,C\}}, \mathcal{G} = \{1, 2, 3\}$

Proof of De Bruijn-Erdős:

Let \mathcal{G} be a set system which satisfies the conditions of the theorem and let \mathcal{F} be its dual.

Case 1: We obtained the same set twice by the construction of \mathcal{F} , so there are $A, B \in \mathcal{F}$, $A=B$, so \mathcal{F} is not a set but a multiset.

In this case A and B contain the same elements. What does it mean for \mathcal{G} ?
 A and B correspond to elements a and b , respectively and any set of \mathcal{G} which contains a or b it contains the other. But there is exactly one such set. If c is another element from $[m]$, then there must be a set which contains a and c , but then it also contains b . So c is also contained in that set. Therefore $\mathcal{G} = \{[m]\}$.

Case 2: Every set appears only once during the construction of \mathcal{F} , so it is a set of subsets.

Then $\forall a, b \in [m] \Rightarrow \exists! C \in \mathcal{G} : a, b \in C \Rightarrow \forall A, B \in \mathcal{F} \Rightarrow \exists! c \in [m] : c \in A \cap B$.

So we can apply Fischer's inequality and we obtain that $m \leq n = |\mathcal{G}|$. \square

m can be attained by the following example which is called as the "near pencil":

