

# FPT suspects and tough customers: Open problems of Downey and Fellows

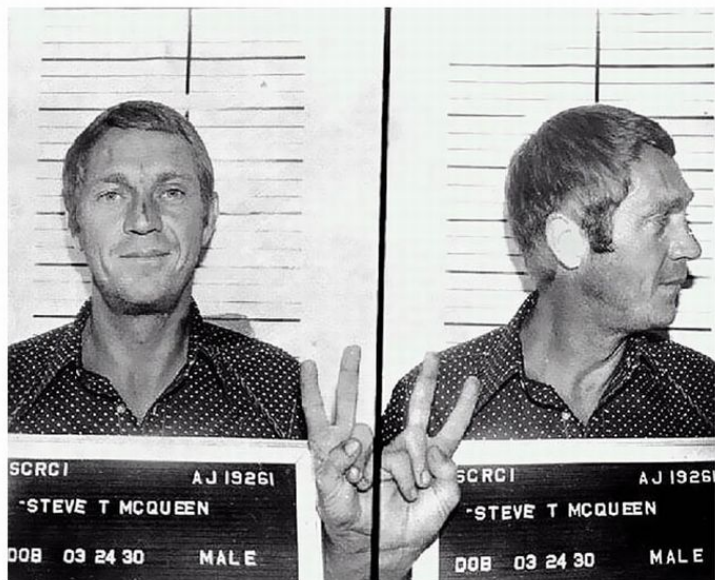
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- We revisit the open problem list of the Downey-Fellows book.
- Good open problems are also significant scientific contributions.
- Were they good problems?
  - Not too easy?
  - Not impossible?
  - Any positive results?

## FPT suspects



# FPT suspects

**Directed Feedback Vertex Set**

*Instance:* A directed graph  $D = (V, A)$ , an integer  $k$ .

*Parameter:*  $k$

*Question:* Is there a set  $S$  of  $k$  vertices such that each directed cycle of  $G$  contains a member of  $S$ ?

DFVS



DFVS

Directed Feedback Vertex Set

**Instance:** A directed graph  $G$

**Parameter:** A positive integer  $k$

**Question:** Is there a set  $S$  of  $k$  vertices such that each directed cycle of  $G$  contains a member of  $S$ ?

- FPT by [Chen et al. 2008]

Topological Containment

**Instance:** An undirected graph  $G$

**Parameter:** A graph  $H$

**Question:** Is  $H$  topologically contained in  $G$ ?

- In XP by [Robertson and Seymour, GM13]
- FPT by [Grohe et al. 2011]

Immersion Order Test

**Instance:** An undirected graph  $G$

**Parameter:** A graph  $H$

**Question:** Does  $H$  has an immersion in  $G$ ?

- FPT by reduction to Topological Containment.

# Open

## Planar Directed Disjoint Paths

**Instance:** A directed planar graph  $G$  and  $k$  pairs  $\langle r_1, s_1 \rangle, \dots, \langle r_k, s_k \rangle$  of vertices of  $G$

**Parameter:**  $k$

**Question:** Does  $G$  have  $k$  vertex-disjoint paths  $P_1, \dots, P_k$  with  $P_i$  running from  $r_i$  to  $s_i$ ?

- In XP by [Schrijver 1994]



Planar  $t$ -Normalized Weighted Satisfiability

**Instance:** A planar  $t$ -normalized formula  $X$

**Parameter:** A positive integer  $k$

**Question:** Does  $X$  have a satisfying assignment of weight  $k$ ?

- What is exactly a planar  $t$ -normalized formula?
- FPT by standard techniques (layering + treewidth arguments or reduction to first order model checking).

## W[1]-hard

Planar Multiway Cut

**Instance:** A weighted undirected planar graph  $G$  with terminals  $\{x_1, \dots, x_k\}$  and a positive integer  $M$

**Parameter:**  $k$

**Question:** Is there a set of edges of total weight  $\leq M$  whose removal disconnects each terminal from all others?

- Can be solved in time  $n^{O(k)}$  by [Dahlhaus et al. 1994].
- Can be solved in time  $2^{O(k)} \cdot n^{O(\sqrt{k})}$  [Klein and M. 2012]
- W[1]-hard and no  $f(k) \cdot n^{o(\sqrt{k})}$  algorithm [M. 2012]

## Tough customers



## W[1]-hard

Fixed Alphabet Longest Common Subsequence (LCS)

**Instance:**  $k$  sequences  $X_i$  over an alphabet  $\Sigma$  of fixed size and a positive integer  $m$

**Parameter:**  $k$

**Question:** Is there a string  $X \in \Sigma^*$  of length  $m$  that is a subsequence of each of the  $X_i$ ?

- $O(n^{k+1})$  time by simple dynamic programming.
- W[1]-hard by [Pietrzak 2003] with binary alphabet.

Crossing Number

**Instance:** An undirected graph  $G$

**Parameter:** A positive integer  $k$

**Question:** Is the crossing number of  $G$  is at most  $k$ ?

- FPT:  $f(k) \cdot n^2$  algorithm by [Grohe 2001]
- $f(k) \cdot n$  algorithm by [Kawarabayashi and Reed 2007]

### Minimum Degree Graph Partition

**Instance:** An undirected graph  $G$

**Parameter:** Positive integers  $k$  and  $d$

**Question:** Can  $V(G)$  be partitioned into disjoint subsets  $V_1, \dots, V_m$  so that for  $1 \leq i \leq m$ ,  $|V_i| \leq k$  and at most  $d$  edges have exactly one endpoint in  $V_i$ ?

- For fixed  $k$  and  $d$ , graphs with such partitions are closed under immersion [Langston and Plaut 1998].
- Immersion is wqo [Robertson and Seymour GM23].
- Immersion testing is FPT [Grohe et al. 2011].
- $\Rightarrow$  Minimum Degree Graph Partition is (nonuniform) FPT.
- $O^*(2^{O(k)})$  and  $O^*(2^{O(d)})$  time by [Lokshtanov and M. 2011].

## Short Cheap Tour

**Instance:** A graph  $G$ , integer  $S$ , and edge weighting  $w: E(G) \rightarrow \mathbb{Z}$

**Parameter:** A positive integer  $k$

**Question:** Is there a tour through at least  $k$  nodes of  $G$  of cost at most  $S$ ?

- *“Using the methods of [PV91] or [AYZ94], it can be shown that the impoverished travelling salesman can visit at least  $k$  cities and return home for a given budget is FPT.”*  
[Fellows 2001]

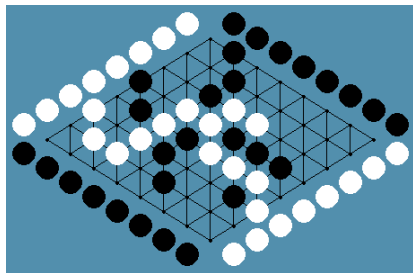
# Open

Short Generalized Hex

**Instance:** An undirected graph  $G$  with two distinguished vertices  $v_1$  and  $v_2$

**Parameter:** A positive integer  $k$

**Question:** Does player one have a winning strategy of at most  $k$  moves in Generalized Hex?





Jump Number

**Instance:** A poset  $P$

**Parameter:** A positive integer  $k$

**Question:** Is the jump number of  $P$  at most  $k$ ?

- In XP by [El-Zahar and Schmerl 1984]
- FPT by [McCartin 2001]

# Open

Shortest Vector

**Instance:** A basis  $X = \{x_1, x_2, \dots, x_t\} \subset \mathbb{Z}^n$  for a lattice  $\mathcal{L}$

**Parameter:** A positive integer  $k$

**Question:** Is there a non-zero vector  $(a_1, \dots, a_n) \in \mathcal{L}$ , such that  $\sum_{i=1}^t a_i^2 \leq k$ ?

- NP-hard under randomized reduction by [Ajtai 1998].
- In XP (trivial).
- Other norms?

# Open

Even Set

**Instance:** An undirected red/blue bipartite graph  $G = (\mathcal{R}, \mathcal{B}, E)$

**Parameter:** A positive integer  $k$

**Question:** Is there a non-empty set of at most  $k$  vertices  $R \subseteq \mathcal{R}$ , such that each member of  $\mathcal{B}$  has an even number of neighbors in  $R$ ?

- Hypergraph formulation.
- Minimum distance of linear codes of  $\text{GF}[2]$ .
- Minimum cycle in a binary matroid.

# Conclusions

- A very good list of problems.
- Only two problems turned out to be  $W[1]$ -hard (one FPT suspect and one tough customer).
- My favorite remaining open problems: Even Set and Planar Directed Disjoint Paths.

Happy birthday  
Mike!