FDBG, the CLP(*FD***) Debugger Library of SICStus Prolog**

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Introduction

CLP(\mathcal{FD}) and **Debugging**

- $CLP(\mathcal{FD})$ implementations:
 - ★ Designated development environments (e.g. OPL Studio);
 - Embedded into a host language Prolog is a logical choice (backtracking, logic variables).
- $CLP(\mathcal{FD})$ and debugging in SICStus Prolog:
 - ★ Extensive $CLP(\mathcal{X})$ libraries, including $CLP(\mathcal{FD})$;
 - ★ An excellent, flexible, extensible debugger for Prolog;
 - * minimal support for $CLP(\mathcal{X})$ debugging (until FDBG).
- Possible approaches to observe a $CLP(\mathcal{FD})$ run:
 - * interactive tools (e.g. step-by-step debuggers);
 - ★ assertion based methods;
 - * trace generation and analysis ideal for nonlinear program execution, like in the case of $CLP(\mathcal{FD})$.

FDBG and its Event Trace

- FDBG = Finite domain DeBuGger
- Main purpose: enable CLP(*FD*) programmers to gather information about constraints and variables possibly even without modifying the observed program.
- FDBG translates the run of a CLP(*FD*) program into an event trace:
 - ★ a sequence of log entries;
 - * each entry corresponds to a $CLP(\mathcal{FD})$ event;
 - ★ an event represents:
 - * the activity of a *constraint* and its effect on variables;
 - * a *labeling* decision while exploring the search tree.
 - ★ appearance of entries is fully customizable.

| ?-

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 fdbg_on.

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    fdvar_1 = inf..sup -> inf..5
    Constraint exited.
```

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<fdvar_1> in 4..sup
fdvar_1 = inf..5 -> 4..5
Constraint exited.
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Labeling [3, <fdvar_1>]: starting in range 4..5.
Labeling [3, <fdvar_1>]: step: <fdvar_1> = 4
X = 4 ? ;
Labeling [3, <fdvar_1>]: step: <fdvar_1> >= 5
X = 5 ? ;
Labeling [3, <fdvar_1>]: failed.
```

Basic Concepts

CLP(*FD*) Events

- $CLP(\mathcal{FD})$ problem solving consists of two repeated phases:
 - narrowing a variable domain due to constraint propagation;
 - ★ narrowing a variable domain due to labeling.
- Observation: with two classes of events we can describe the behavior of a CLP(*FD*) program:
 - ★ constraint events
 - * a constraint is woken up and performs propagation;
 - ★ labeling events
 - * a choicepoint is created or exhausted (through failure);
 - * the domain of a variable is narrowed.
- Events are intercepted and dispatched to *visualizers* by the FDBG core.

Visualizers

- Predicates responsible for handling $CLP(\mathcal{FD})$ events;
- Usually display trace information;
- In general can do any kind of processing (like checking invariants);
- Analogously to event classes, there are two types:
 - ★ constraint visualizers;
 - \star labeling visualizers.
- FDBG provides default built-in visualizers for both types;
- Utility predicates support writing custom visualizers.

The User Interface

An Example – The N-queens Problem

```
nqueens(N, Queens) :-
        bb_put(board_size, N),
        length(Queens, N),
        fdbg_assign_name(Queens, queen),
        domain(Queens, 1, N),
        constrain_all(Queens),
                                          % break symmetry
%
        asiymmetric(N, Queens),
        labeling([ff], Queens).
no_threat(X, Y, I) :-
        fd_global(no_threat(X,Y,I), 1, [val(X),val(Y)]).
:- multifile clpfd:dispatch_global/4.
clpfd:dispatch_global(no_threat(X,Y,I), S, S, Actions) :-
            integer(X) -> no_threat_prop(Y, X, I, Actions)
        (
            integer(Y) -> no_threat_prop(X, Y, I, Actions)
        ;
                           Actions = []
        ;
        ).
. . .
```

Built-in Visualizers

- FDBG uses built-in visualizers by default;
- Built-in visualizers can work without any program modification.

One block of output of the constraint visualizer

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no_threat(2,<queen_3>,2)
  queen_3: 1..2 -> {1}
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- The legend lists all the variables of the constraint;
- Most common behaviors are entailment (above) and failure.
- Variable names:
 - ★ all variables are assigned a name for clarity;
 - ★ needed because name in source is not preserved in Prolog;
 - ★ usually displayed between angle brackets (<queen_3>).

Output of the labeling visualizer

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Details:

• Each event results in one line of output;

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 - ★ failure.

Variable Naming

- Variables are assigned a name:
 - * manually by calling fdbg_assign_name/2;
 - * automatically when calling fdbg_annotate/2,3;
 - \star auto-assigned names look like fdvar_N (N unique counter).
- Names are primarily used to refer to variables in the trace;
- This done via *annotation:* each variable in a term is replaced by a term containing its name, itself, and its narrowed domain;
- Convenience service: assign names to an entire term and each variable in it with a single call:

term/variable	selector	name	
bar(A, [B, C])	[]	foo	assigned name

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term/variable	selector	name	
bar(A, [B, C])	[]	foo	assigned name
Α	[1]	foo_1) implicit
В		foo_2_1	<pre>implicit derived names</pre>
C	[2,#2]	foo_2_2	

Customizing Visualizers

- Built-in visualizers have no knowledge of the problem structure;
- Customized visualization can exploit this additional knowledge;
- Customization is possible on two levels:
 - slight modification of the output of the built-in visualizers by defining hook predicates;
 - ★ writing custom visualizers.
- Use a hook predicate to modify the legend of *N*-queens:

```
* no_threat(2,<queen_3>,2)
    queen_3: [ X - . . ]
    Constraint exited.
```

• Or a custom visualizer to completely redefine it:

```
* no_threat(4,<queen_3>,1)
    [ X X . . ]
    [ . . . X ]
    [ X X - - ]
    [ X X X X ]
```

Legend Portray Hook

```
fdbg:legend_portray(Name, Var, After) :-
        write(Name), write(': '),
        print_row(Var, After).
print_row(Var, After) :-
        bb_get(board_size, N),
        fd_set(Var, Now),
        write('['), print_fields(1, N, Now, After), write(']').
print_fields(I, N, _, _) :-
        I > N, !, write(', ').
print_fields(I, N, Now, After) :-
        write(' ').
        ( fdset_member(I, After) -> write('X') % allowed
           fdset_member(I, Now) -> write('-') % being pruned
                                      write('.') % pruned
        ,
        ),
        I1 is I+1,
        print_fields(I1, N, Now, After).
```

Custom Visualizer

```
nqueens_show(Constraint, Actions) :-
        fdbg_current_name(Queens, queen),
        fdbg_annotate(Constraint, AConst, _),
        fdbg_annotate(Queens, Actions, AQueens, _),
        print(AConst), nl,
        print_board(AQueens).
print_board([]) :- nl.
print_board([fdvar(_,Var,After)|Qs]) :- !,
        write('').
        print_row(Var, After), nl,
        print_board(Qs).
print_board([V|Qs]) :-
        write(' '),
        fdset_singleton(Set, V),
        print_row(V, Set), nl,
        print_board(Qs).
```

Implementation Issues

Event detection

- SICStus debugger provides *advice points*: programmable breakpoints;
- FDBG places advice points on all constraint handling predicates;
- Limitation: only *global constraints* are handled, *indexicals* are ignored;
- Workaround: when FDBG is turned on, constraints otherwise compiled as indexicals translate into global constraints (through goal expansion);
- Consequences:
 - ⋆ FDBG should be consulted before the program to be traced;
 - Negligible effect on performace compared to overhead of FDBG in general;
 - Minor behavioral changes (slightly different propagation);
 - ★ Original form of constraints is lost in the process.

Variable Naming Revisited

Reminder

- Names can be assigned by the user to any term or variable;
- Visualizers refer to variables with names exploiting *annotation*;
- *Annotation* is the process of replacing variables in a term with descriptive compound terms.

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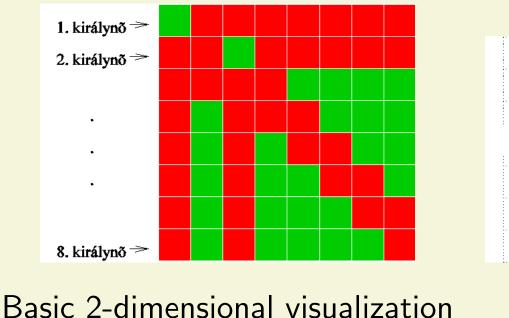
Implementation

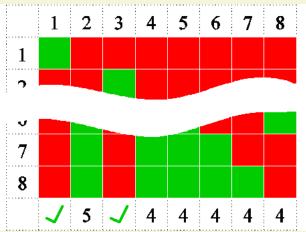
- Names are stored in an AVL tree in the *global private* field;
- As a result, the name store is *volatile* (fresh for each query);
- Consequences:
 - ★ No need to clear the name store after each query (good);
 - Need to assign names in each query (seemingly inconvenient but unavoidable: different variables!);
 - \star The best place to do this is within the program itself.

Future Work and Conclusions

Future Work

- Indexical to global constraint redirection fails to preserve original form of constraints – need to find an acceptible solution;
- When labeling fails, there is no information in the log about what state the domains of variables are *restored* to through backtracking;
- A generic, configurable graphical visualizer plans have already been proposed, still need to implement it.





Embellished version

Conclusions

- Presented basic trasing scheme for CLP(*FD*) programs written in SICStus Prolog;
- Introduced events and visualizers;
- Showed how variable naming and output customization can help to clarify the trace log;
- Given examples to visualizer customization;
- Covered a few implementational details;
- Scetched possible directions of future development.

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Thanks for your attention!